

# Lecture 5: Shor's algorithm

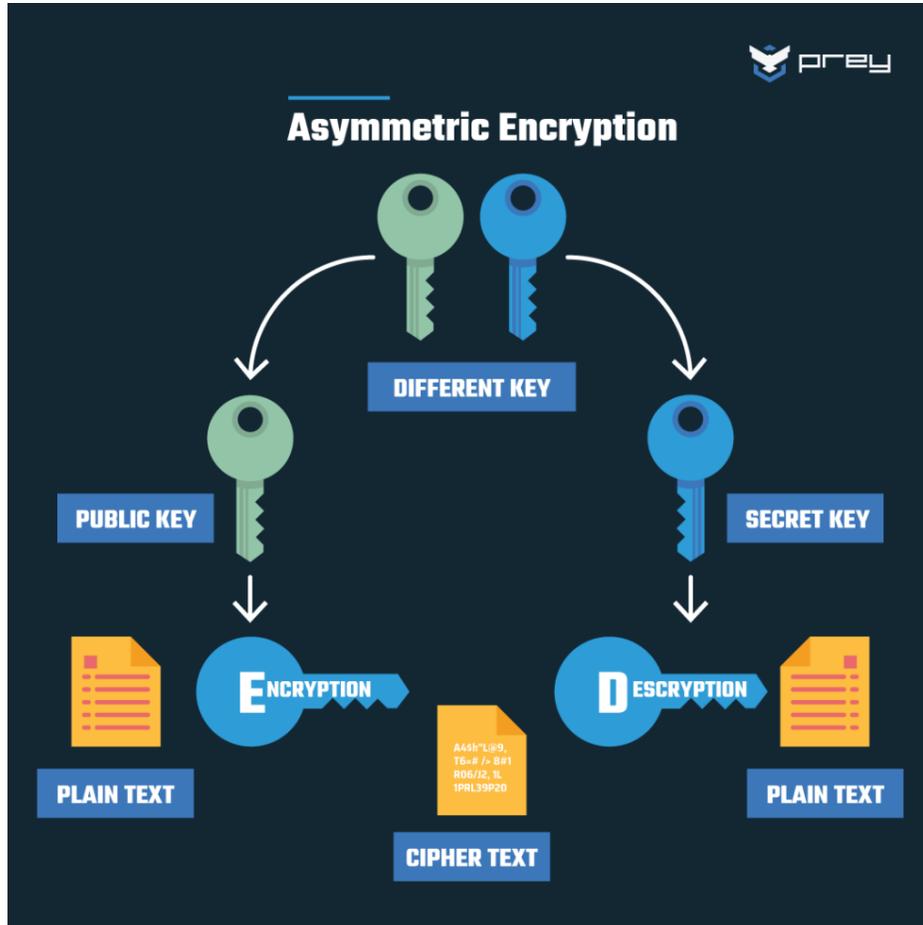
January 28, 2026



# Known algorithms with exponential speedup

| Name                   | $G$  | $X$  | $K$   | Function   |
|------------------------|--|--|---|--|
| Deutsch                | $\{0, 1\}, \oplus$   | $\{0, 1\}$                                     | $\{0\}$ or $\{0, 1\}$                           | $K = \{0, 1\} : \begin{cases} f(x) = 0 \\ f(x) = 1 \end{cases}$ $K = \{0\} : \begin{cases} f(x) = x \\ f(x) = 1 - x \end{cases}$ |
| Simon                  | $\{0, 1\}^n, \oplus$   | any finite set                                 | $\{0, s\}$<br>$s \in \{0, 1\}^n$                | $f(x \oplus s) = f(x)$   |
| Period-finding         | $\mathbf{Z}, +$  | any finite set                                 | $\{0, r, 2r, \dots\}$<br>$r \in G$              | $f(x + r) = f(x)$  |
| Order-finding          | $\mathbf{Z}, +$  | $\{a^j\}$<br>$j \in \mathbf{Z}_r$<br>$a^r = 1$ | $\{0, r, 2r, \dots\}$<br>$r \in G$              | $f(x) = a^x$<br>$f(x + r) = f(x)$  |
| Discrete logarithm     | $\mathbf{Z}_r \times \mathbf{Z}_r$<br>$+ \pmod{r}$           | $\{a^j\}$<br>$j \in \mathbf{Z}_r$<br>$a^r = 1$ | $(\ell, -\ell s)$<br>$\ell, s \in \mathbf{Z}_r$ | $f(x_1, x_2) = a^{kx_1 + x_2}$<br>$f(x_1 + \ell, x_2 - \ell s) = f(x_1, x_2)$  |
| Order of a permutation | $\mathbf{Z}_{2^m} \times \mathbf{Z}_{2^n}$<br>$+ \pmod{2^m}$ | $\mathbf{Z}_{2^n}$                             | $\{0, r, 2r, \dots\}$<br>$r \in X$              | $f(x, y) = \pi^x(y)$<br>$f(x + r, y) = f(x, y)$<br>$\pi = \text{permutation on } X$  |
| Hidden linear function | $\mathbf{Z} \times \mathbf{Z}, +$                            | $\mathbf{Z}_N$                                 | $(\ell, -\ell s)$<br>$\ell, s \in X$            | $f(x_1, x_2) =$<br>$\pi(sx_1 + x_2 \pmod{N})$<br>$\pi = \text{permutation on } X$  |
| Abelian stabilizer     | $(H, X)$<br>$H = \text{any Abelian group}$                   | any finite set                                 | $\{s \in H \mid f(s, x) = x, \forall x \in X\}$ | $f(gh, x) = f(g, f(h, x))$<br>$f(gs, x) = f(g, x)$   |

# Cryptography using RSA



## RSA Algorithm

### Key Generation

|                                  |   |
|----------------------------------|---|
| Select $p, q$                    | $p$ and $q$ , both prime; $p \neq q$    |
| Calculate $n = p \times q$       |   |
| Calculate $\phi(n) = (p-1)(q-1)$ |   |
| Select integer $e$               | $\gcd(\phi(n), e) = 1; 1 < e < \phi(n)$ |
| Calculate $d$                    | $de \bmod \phi(n) = 1$                  |
| Public key                       | $KU = \{e, n\}$                         |
| Private key                      | $KR = \{d, n\}$                         |

### Encryption

|             |                   |
|-------------|-------------------|
| Plaintext:  | $M < n$           |
| Ciphertext: | $C = M^e \pmod n$ |

### Decryption

|             |                   |
|-------------|-------------------|
| Plaintext:  | $M$               |
| Ciphertext: | $M = C^d \pmod n$ |

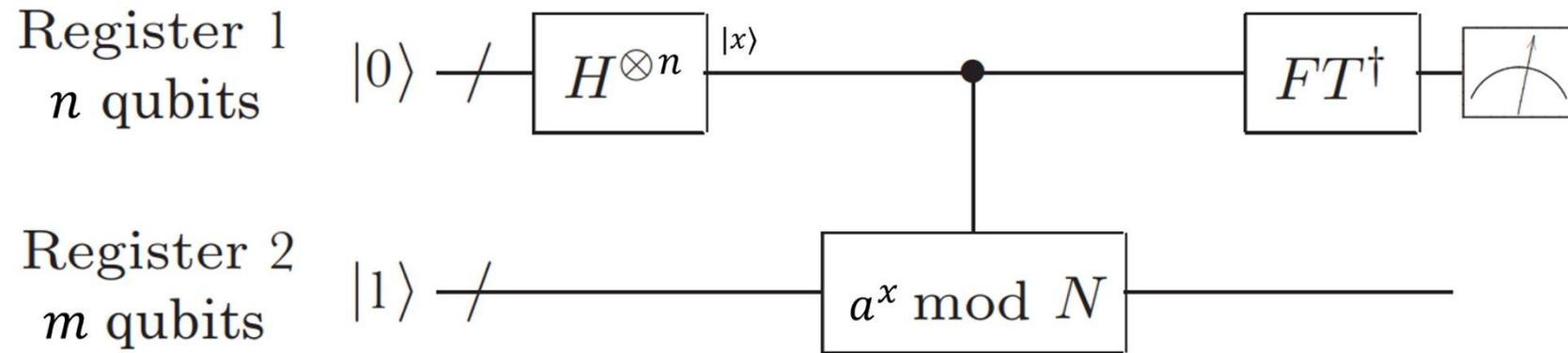
How to factor 2048 bit RSA integers with less than a million noisy qubits

Craig Gidney

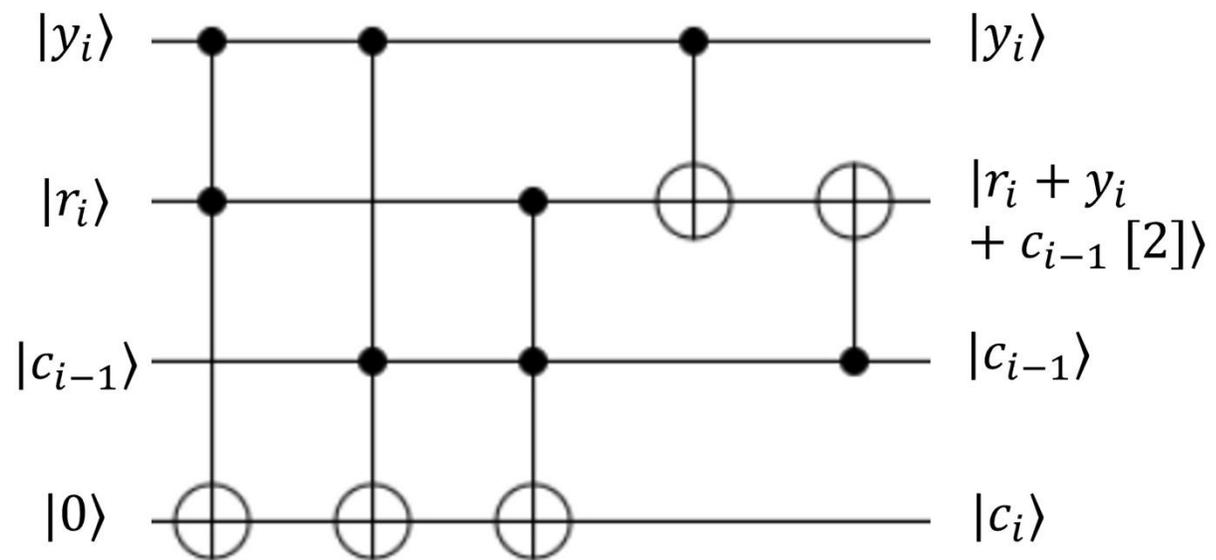
Google Quantum AI, Santa Barbara, California 93117, USA

June 9, 2025

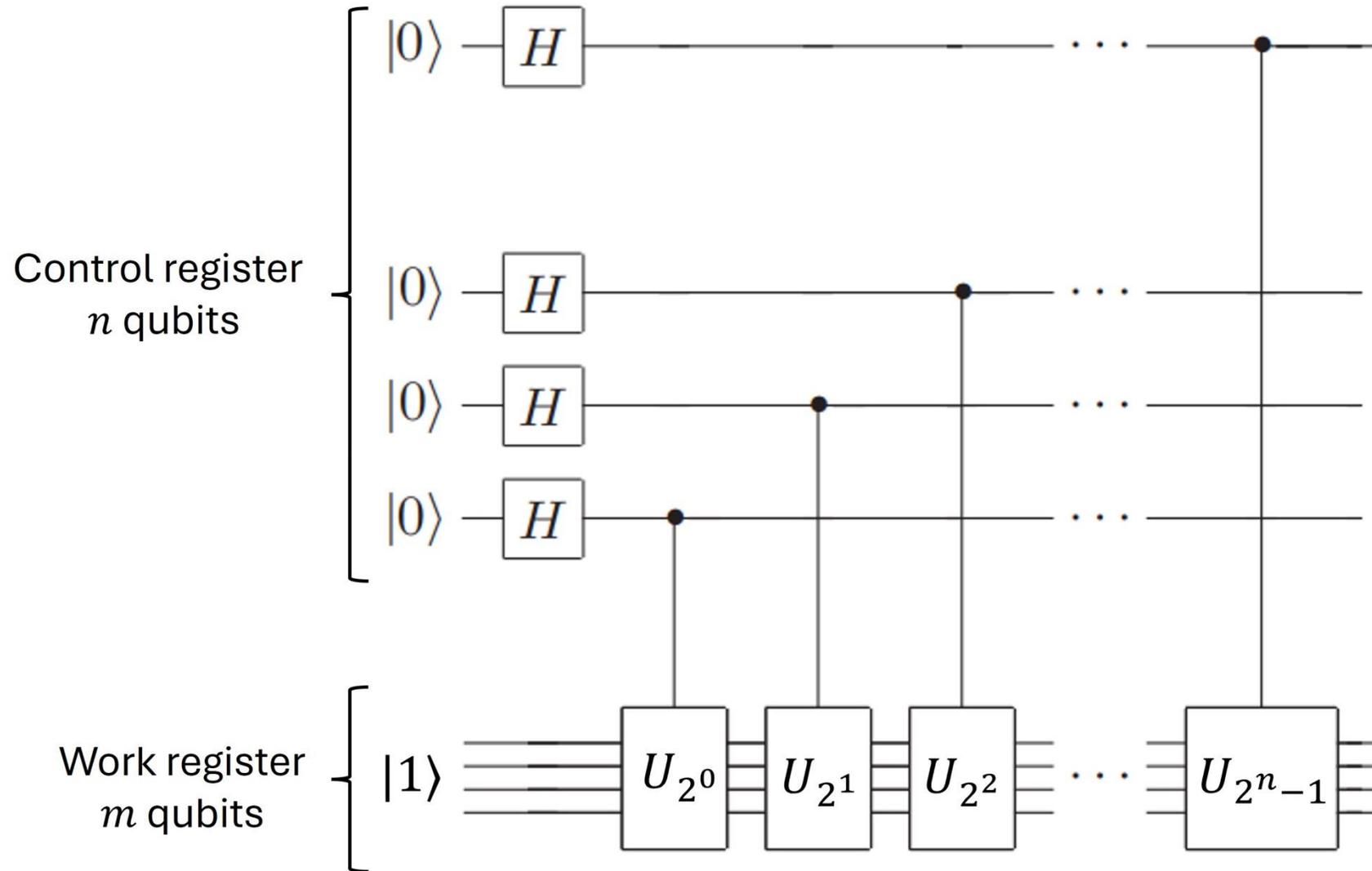
# Quantum part of Shor



# Adder



# Modular exponentiation



# QFT

